## Title

## Decomposition of nonnegative singular matrices into product of nonnegative idempotent matrices

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Decomposition of nonnegative singular matrices into product of nonnegative idempotent matrices and mORE...(skew) codes.

ALGEBRAIC METHODS IN CODING THEORY
Ubatuba July 2017

## Pioneers

- J.M.Howie (1966) The maps from a finite set to itself that are not onto can be presented as products of idempotents.
- J.A. Erdos (1968): singular matrices over fields.
- J. Laffey (1983): singular matrices over commutative euclidean domains.
- Hannah and O'Meara decomposition of some elements in a von Neumann ring into product of idempotent elements.
- Bhaskara Rao (2009) considered matrices over commutative PID's.
- W. Ruitenberg (1993) Matrices over Hermite domains.
- There are connections between decompositions into products of idempotents and factorizations of invertible matrices into product of elementary matrices. (Facchini-Leroy(2016), Salce-Zanardo,...)


## Examples and particular decompositions

## Examples

$$
\text { (a) }\left(\begin{array}{ll}
a & 0 \\
0 & 0
\end{array}\right)=\left(\begin{array}{ll}
1 & 1 \\
0 & 0
\end{array}\right)\left(\begin{array}{cc}
1 & 0 \\
a-1 & 0
\end{array}\right) \text {. }
$$

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$$
\begin{aligned}
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a & 0 \\
0 & 0
\end{array}\right) & =\left(\begin{array}{ll}
1 & 1 \\
0 & 0
\end{array}\right)\left(\begin{array}{cc}
1 & 0 \\
a-1 & 0
\end{array}\right) . \\
\left(a^{\prime}\right)\left(\begin{array}{ll}
a & 0 \\
0 & 0
\end{array}\right) & =\left(\begin{array}{ll}
1 & a \\
0 & 0
\end{array}\right)\left(\begin{array}{ll}
0 & 0 \\
0 & 1
\end{array}\right)\left(\begin{array}{ll}
1 & 0 \\
1 & 0
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\end{aligned}
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(a) $\left(\begin{array}{ll}a & 0 \\ 0 & 0\end{array}\right)=\left(\begin{array}{ll}1 & 1 \\ 0 & 0\end{array}\right)\left(\begin{array}{cc}1 & 0 \\ a-1 & 0\end{array}\right)$.
(a') $\left(\begin{array}{ll}a & 0 \\ 0 & 0\end{array}\right)=\left(\begin{array}{ll}1 & a \\ 0 & 0\end{array}\right)\left(\begin{array}{ll}0 & 0 \\ 0 & 1\end{array}\right)\left(\begin{array}{ll}1 & 0 \\ 1 & 0\end{array}\right)$.
(b) $\left(\begin{array}{cc}a & a c \\ 0 & 0\end{array}\right)=\left(\begin{array}{cc}1 & 1+c \\ 0 & 0\end{array}\right)\left(\begin{array}{cc}1-c a+c & c-c a c+c^{2} \\ a-1 & a c-c\end{array}\right)$.

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(a') $\left(\begin{array}{ll}a & 0 \\ 0 & 0\end{array}\right)=\left(\begin{array}{ll}1 & a \\ 0 & 0\end{array}\right)\left(\begin{array}{ll}0 & 0 \\ 0 & 1\end{array}\right)\left(\begin{array}{ll}1 & 0 \\ 1 & 0\end{array}\right)$.
(b) $\left(\begin{array}{cc}a & a c \\ 0 & 0\end{array}\right)=\left(\begin{array}{cc}1 & 1+c \\ 0 & 0\end{array}\right)\left(\begin{array}{cc}1-c a+c & c-c a c+c^{2} \\ a-1 & a c-c\end{array}\right)$.
(c) $\left(\begin{array}{cc}a c & a \\ 0 & 0\end{array}\right)=\left(\begin{array}{ll}1 & a \\ 0 & 0\end{array}\right)\left(\begin{array}{ll}0 & 0 \\ c & 1\end{array}\right)$,

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(a') $\left(\begin{array}{ll}a & 0 \\ 0 & 0\end{array}\right)=\left(\begin{array}{ll}1 & a \\ 0 & 0\end{array}\right)\left(\begin{array}{ll}0 & 0 \\ 0 & 1\end{array}\right)\left(\begin{array}{ll}1 & 0 \\ 1 & 0\end{array}\right)$.
(b) $\left(\begin{array}{cc}a & a c \\ 0 & 0\end{array}\right)=\left(\begin{array}{cc}1 & 1+c \\ 0 & 0\end{array}\right)\left(\begin{array}{cc}1-c a+c & c-c a c+c^{2} \\ a-1 & a c-c\end{array}\right)$.
(c) $\left(\begin{array}{cc}a c & a \\ 0 & 0\end{array}\right)=\left(\begin{array}{ll}1 & a \\ 0 & 0\end{array}\right)\left(\begin{array}{ll}0 & 0 \\ c & 1\end{array}\right)$,
(d) with $b \in U(R),\left(\begin{array}{ll}a & b \\ 0 & 0\end{array}\right)=\left(\begin{array}{cc}b\left(b^{-1} a\right) & b \\ 0 & 0\end{array}\right)$ is factorized as in (c).

## Particular matrices

## Theorem

(a) If $R$ is a ring and $A \in M_{n}(R)$ is strictly upper triangular then $A$ is a product of idempotent matrices.
(b) If $n>1$ and a matrix $A \in M_{n}(\mathbb{R})$ (resp. $A \in M_{n}\left(\mathbb{R}^{+}\right)$) has only one nonzero row, then it is a product of (resp. nonnegative) idempotent matrices.

## Question and particular matrices

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Main Question : Can real nonnegative singular matrices be decomposed into product of nonnegative idempotents ?

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## Lemma (Particular matrices)

(a) If $B \in M_{n \times n}\left(\mathbb{R}^{+}\right)$is an $n \times n$ matrix which is a product of nonnegative idempotents, then the same is true for the matrix $\left(\begin{array}{ll}B & C \\ 0 & 0\end{array}\right)$ where $C \in M_{n \times 1}(\mathbb{R})\left(\right.$ resp. $\left.C \in M_{n \times 1}\left(\mathbb{R}^{+}\right)\right)$and the other blocks are of appropriate sizes.
(b) If $A \in M_{n}(\mathbb{R})\left(\right.$ resp. $A \in M_{n}\left(\mathbb{R}^{+}\right)$), $n \geq 3$, has all its $i^{t h}$ rows and columns zero whenever $i \geq 3$, then $A$ is a product of (resp. nonnegative ) idempotent matrices.

## Rank one

## Proposition

Let $A \in M_{n}\left(\mathbb{R}^{+}\right), n>1$, be a nonnegative matrix of rank 1 . Then $A$ is a product of nonnegative idempotent matrices.

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## Remark (A.Alahamadi,S.K. Jain, A.L., Sathaye,2016)

It can be shown that in fact rank 1 nonnegative matrices can be decomposed into a product of three idempotent nonnegative matrices.

## Rank two

## Theorem

Let $A \in M_{n}\left(\mathbb{R}^{+}\right), n>2$, be a nonnegative singular matrix of rank 2. Then $A$ is a product of nonnegative idempotent matrices.

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The proof is based on the following easy lemma:

## Lemma

Let $S \subset\left(\mathbb{R}^{+}\right)^{n}$ be a finite set such that $\operatorname{dim}_{\mathbb{R}}<S>\leq 2$. Then there exist $s_{1}, s_{2} \in S$ such that every element of $S$ is a positive linear combination of $s_{1}$ and $s_{2}$.

## Counter-example

For singular nonnegative matrices of higher rank the decomposition does not necessarily exist:

## Example

$$
A_{\alpha}:=\left(\begin{array}{cccc}
\alpha & \alpha & 0 & 0 \\
0 & 0 & \alpha & \alpha \\
\alpha & 0 & \alpha & 0 \\
0 & \alpha & 0 & \alpha
\end{array}\right), \quad \text { where } \alpha \in \mathbb{R}^{+}, \alpha \neq 0
$$

If $A_{\alpha}=E_{1} \ldots E_{n}$ is such that $E_{i}^{2}=E_{i} \in M_{n}\left(\mathbb{R}^{+}\right)$then $A_{\alpha}=A_{\alpha} E_{n}$ and a direct computation shows that $E_{n}=I d$.. Remark that $A_{\frac{1}{2}}$ is a positive doubly stochastic matrix.

## Nilpotent matrices

## Proposition

If $A$ is Nonnegative nilpotent there exists a permutation matrix such that $P A P^{t}$ is an upper triangular nonnegative matrix.

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## Corollary

Nonnegative nilpotent matrices are product of nonnegative idempotent matrices.

## quasi-permutation matrices

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A matrix $A \in M_{n, n}\left(\mathbb{R}^{+}\right)$is a quasi-permutation matrix if each row and each column has at most one nonzero element.

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## Theorem

A nonnegative singular quasi-permutation matrix is always a product of nonnegative idempotent matrices.

## Nonnegative Von Neumann inverses

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For a nonnegative matrix $A$ to have a nonegative von Neumann inverse, it must be of a very special form (quasi permutation by block with all blocks of rank one). Using this form and the previous result on quasi-permutation matrices we get the following theorem.

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## Theorem

A nonnnegative singular matrix with nonnegative von Neumann inverse is a product of nonnegative idempotent matrices.

## Periodic matrices

## Definitions

(1) A matrix $A$ is periodic if there exist positive integers $l<s$, such that $A^{\prime}=A^{s}$

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## Theorem

Let $A$ be a nonnegative periodic matrix with no zero row or zero column. If either the index of $A$ is 1 or $A>A^{n}$ for some $n$, then $A$ is a product of nonnegative idempotent matrices.

## 0-1 definite matrices

## Definition

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## Theorem

Let $A \in M_{n}(\mathbb{R})$ be a singular definite $0-1$ matrix. Then $A$ is a product of nonnegative idempotent matrices.

## Plan

(1) A) Ore extensions.
(2) B) Polynomial maps.
(3) C) Pseudo-linear transformations.
(1) D) $(\sigma, \delta)$-codes.
(6) E) $\mathrm{W}(\sigma, \delta)$-codes.

## Ore Extension

B) Polynomial maps and roots
C) Pseudo linear transformations $(\sigma, \delta)$-codes

## Layout

(1) History
(2) Particular decompositions
(3) Nonnegative singular matrices

- special families of nonnegative matrices
(5) ...and mORE
- Ore Extension
- B) Polynomial maps and roots
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C) Pseudo linear transformations $(\sigma, \delta)$-codes


## Ore extensions

$A$ a ring, $D$ a derivation of $A$, for $a \in A L_{a}$ is the left multiplication by $a$.

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D \circ L_{a}=L_{a} \circ D+L_{D(a)}
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More generally:
Assume the polynomials in $X$ can be written as $\sum a_{i} X^{i}$ and there are maps $\sigma, \delta$ from $A$ to $A$ such that

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Then associativity of the product will give that $\sigma \in \operatorname{End}(A)$ and $\delta$ is a $\sigma$ derivation i.e. $\delta \in \operatorname{End}(A,+)$ and

$$
\delta(a b)=\sigma(a) \delta(b)+\delta(a) b
$$

C) Pseudo linear transformations $(\sigma, \delta)$-codes

## Examples

The set of these polynomials form a ring denoted by
$R=A[X ; \sigma, \delta]$ (O. Ore, 1930's)
(1) $R=\mathbb{C}[t ;-]$; we have $t i=-i t$ and $t^{2} a=t(\bar{a} t)=a t^{2}$. $\frac{R}{R\left(t^{2}+1\right)} \cong \mathbb{H}$.

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(2) $p$ a prime, $q=p^{\prime}$ and $R=\mathbb{F}_{q}[t ; \sigma]$; where $\sigma(x)=x^{p}$. The center of $R$ is $\mathbb{F}_{p}\left[t^{\prime}\right]$.

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(2) $p$ a prime, $q=p^{\prime}$ and $R=\mathbb{F}_{q}[t ; \sigma]$; where $\sigma(x)=x^{p}$. The center of $R$ is $\mathbb{F}_{p}\left[t^{\prime}\right]$.
(3) $k$ a field, $A_{1}=k[x]\left[y ; I d ., \frac{d}{d x}\right]$ the first Weyl algebra.

- If $\operatorname{char}(k)=p>0, Z\left(A_{1}\right)=k\left[x^{p}, y^{p}\right]$
- If $\operatorname{char}(k)=0$ then $Z\left(A_{1}\right)=k$ and $A_{1}$ is simple.
C) Pseudo linear transformations $(\sigma, \delta)$-codes


## Inner and not inner

The $\sigma$ inner derivation induced by an element $a \in A$ is defined by $\delta_{a} \in \operatorname{End}(A,+)$ by $\delta_{a}(x)=a x-\sigma(x)$ a, for $x \in A$.

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Such a derivation can be "erased": $A\left[t, \sigma, \delta_{a}\right]=A[t-a, \sigma]$. Finite ring can have non inner $\sigma$-derivation even if $\sigma \neq I d$..

## Example

Let $q=p^{\prime}, p$ a prime and $A$ be the subring of $M_{2}\left(\mathbb{F}_{q}\right)$ given by

$$
A=\left\{\left.\left(\begin{array}{ll}
a & b \\
0 & c
\end{array}\right) \right\rvert\, a, b \in \mathbb{F}_{q}, c \in \mathbb{F}_{p}\right\} .
$$

Define $\sigma$ and $\delta$ as follows:

$$
\sigma\left(\left(\begin{array}{ll}
a & b \\
0 & c
\end{array}\right)\right)=\left(\begin{array}{cc}
a^{p} & b^{p} \\
0 & c
\end{array}\right) \quad \text { and } \quad \delta\left(\left(\begin{array}{ll}
a & b \\
0 & c
\end{array}\right)\right)=\left(\begin{array}{cc}
0 & b^{p} \\
0 & 0
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Ore Extension
B) Polynomial maps and roots
C) Pseudo linear transformations
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## Layout

(1) History
(2) Particular decompositions
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(4) special families of nonnegative matrices
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## Polynomial maps

$f(t) \in R=A[t ; \sigma, \delta], a \in A$, there exist $q(t) \in R$ such that $f(t)-q(t)(t-a) \in A$. This element is naturally defined to be the evaluation of $f(t)$ at $a$, denoted $f(a)$.

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f(t)=q(t)(t-a)+f(a)
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Let us compute: $t^{2}=t(t-a)+t a=t(t-a)+\sigma(a) t+\delta(a)=$ $t(t-a)+\sigma(a)(t-a)+\sigma(a) a+\delta(a)$

Hence $\quad t^{2}(a)=\sigma(a) a+\delta(a)$
We will write $N_{i}(a)$ instead of $t^{i}(a)$.

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We will write $N_{i}(a)$ instead of $t^{i}(a)$. Exercise: Compute $N_{3}(a)$ Recurrence formulas:

$$
N_{0}(a)=1, \quad N_{1}(a)=a, \quad N_{i+1}(a)=\sigma\left(N_{i}(a)\right) a+\delta\left(N_{i}(a)\right)
$$

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## Roots

For $f(t)=\sum_{i=0}^{n} a_{i} t^{i} \in R$ and $a \in A$ we have $f(a)=\sum_{i=0}^{n} a_{i} N_{i}(a) . a \in A$ is a right root of $f(t)$ if $f(a)=0$.

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## Examples

(1) If $\sigma=I d$. and $\delta=0$ we have the usual evaluation $N_{i}(a)=a^{i}$. But $A$ can be non commutative so $j$ is not a right root of $(x-j)(x-i) \in \mathbb{H}[x]$.
(2) Many (right) roots: $f(x)=x^{2}+1 \in \mathbb{H}[x]$ then $f\left(\right.$ yiy $\left.^{-1}\right)=0$ for $0 \neq y \in \mathbb{H}$.
(3) (Wedderburn) $D$ a division ring $f(x) \in Z(D)[x]$ and $d \in D$ such that $f(d)=0$ then there exists elements $a_{1}, \ldots a_{n} \in D \backslash 0$ such that

$$
f(x)=\left(x-d^{a_{1}}\right) \ldots\left(x-d^{a_{n}}\right) .
$$

## Examples

## More examples:

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(1) Consider $t^{2} \in A_{1}(k)=k[x]\left[t\right.$; ld., $\left.\frac{d}{d x}\right]$ we have $t^{2}=\left(t-\frac{1}{x}\right)\left(t+\frac{1}{x}\right)$.
(2) Gordon Motzkin: Let $D$ be a division ring and $f(x) \in D[x]$ the roots of $F(x)$ in $D$ belong to at most $\operatorname{deg}(f)$ conjugacy classes

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A nice formula: let $f(t), g(t) \in R=D[t ; \sigma, \delta]$ where $D$ is a dvision ring and $a \in D$.

$$
(f g)(a)=\left\{\begin{aligned}
0 & \text { if } g(a)=0 \\
f\left(a^{g(a)}\right) g(a) & \text { if } g(a) \neq 0
\end{aligned}\right.
$$

where for $a \in D$ and $c \in D^{*}$ we define $a^{c}=\sigma(c) a c^{-1}+\delta(c) c^{-1}$
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## Pseudo linear transformations

$A, \sigma, \delta$, as usual $R=A[t ; \sigma, \delta]$

## Definition

Let ${ }_{A} V$ be a left module. A map $T \in \operatorname{End}(V,+)$ is a P.L.T. if

$$
T(\alpha v)=\sigma(\alpha) T(v)+\delta(\alpha) v \quad \forall \alpha \in A, \forall v \in V
$$

${ }_{A} V$ then becomes a left $R$-module: $\left(\sum_{i=0}^{n} a_{i} t^{i}\right) \cdot v=\sum_{i=0}^{n} a_{i} T^{i}(v)$ for $v \in V$

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Left $R$-modules $\Leftrightarrow$ Left $A$-module and a P.L.T.

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$$

${ }_{A} V$ then becomes a left $R$-module: $\left(\sum_{i=0}^{n} a_{i} t^{i}\right) \cdot v=\sum_{i=0}^{n} a_{i} T^{i}(v)$ for $v \in V$

Left $R$-modules $\Leftrightarrow$ Left $A$-module and a P.L.T.

## Examples

(1) $\delta$ is a P.L.T. defined on $V=A$
(2) Let $C \in M_{n}(A)$ then $T_{C}: A^{n} \longrightarrow A^{n}$ defined by $T_{C}(v)=\sigma(v) C+\delta(v)$ for any $v \in A^{n}$, is a P.L.T.
Andre Leroy (Joint work with A. Alahmadi and S.K. Jain) Product of nonnegative idempotents

## More PLT

## Proposition

Let $R=A[t ; \sigma, \delta]$.
(1) $p(t) \in R, a \in A, p(a)=p\left(T_{a}\right)(1)$
(2) For $x \in U(R) p\left(a^{x}\right) x=p\left(T_{a}\right)(x)$
(3) If $T:{ }_{A} V \longrightarrow_{A} V$ is a PLT, then the map

$$
\phi_{T}: R \longrightarrow \operatorname{End}(V,+): f(t) \mapsto f(T)
$$

is a ring homomorphism.
(4) for $f, g \in R$ and $a \in A$, we have $(f g)(a)=f\left(T_{a}\right)(g(a))$
(0) If $A=D$ is a division ring and $a \in D$ then $\operatorname{ker}\left(P\left(T_{a}\right)\right)$ is a right vector space over the division ring
$C(a):=\left\{x \in D^{*} \mid a^{x}=a\right\} \cup\{0\}$

## Factorizations

## Theorem \{Lam, L.\}

Let $D$ be a division ring, $\sigma \in \operatorname{End}(D)$ and $\delta$ a $\sigma$-derivation. A polynomial $f(t) \in D[t ; \sigma, \delta]$ has roots in I $(\sigma, \delta)$-conjugacy classes $\Delta\left(a_{i}\right):=\left\{a_{i}^{x}=\sigma(x) a_{i} x^{-1}+\delta(x) x^{-1} \mid x \in D^{*}\right\}$. We have

$$
\sum_{i=1}^{l} \operatorname{Dim}_{C_{i}} \operatorname{Ker}\left(f\left(T_{a_{i}}\right)\right) \leq \operatorname{deg}(f(t))
$$

The equality occurs if an only if $f(t)$ is a Wedderburn polynomial.
C) Pseudo linear transformations $(\sigma, \delta)$-codes

## Layout

(1) History

- Particular decompositions
(3) Nonnegative singular matrices
- special families of nonnegative matrices
(5) ...and mORE
- Ore Extension
- B) Polynomial maps and roots
- C) Pseudo linear transformations
- $(\sigma, \delta)$-codes


## Ulmer, Boucher

Just to give an idea: there are 603 different nontrivial right divisors of $t^{14}-1 \in \mathbb{F}_{4}[t ; \theta]$ with $\theta(z)=z^{2}$ comparing with 25 different factors of $x^{14}-1 \in \mathbb{F}_{4}[x]$.
F. Ulmer, D. Boucher started to use skew polynomial rings $(\delta=0)$ to create codes and study them. As an alphabet they not only used fields but also cyclic modules of the form $\frac{R}{R f(t)}$ where $R=F[t ; \sigma]$.

## Example

In $\mathbb{F}_{4}[t ; \theta]$ with $\theta(z)=z^{2}$ where $\alpha \in \mathbb{F}_{4}$ satisfies $\alpha^{2}+\alpha+1=0$, we have: $t^{4}+t^{2}+1=\left(t^{2}+t+1\right)^{2}=\left(t^{2}+\alpha^{2}\right)\left(t^{2}+\alpha\right)=$ $\left(t^{2}+\alpha\right)\left(t^{2}+\alpha^{2}\right)=\left(t^{2}+\alpha^{2} t+1\right)^{2}=\left(t^{2}+\alpha t+1\right)^{2}$,

## $C<\frac{R}{R f}$ with $R=A[t ; \sigma, \delta]$

## Definition

Let $f(t), g(t) \in R=A[t ; \sigma, \delta]$ monic and such that $f(t) \in \operatorname{Rg}(t)$. A subset of $C \subseteq A^{n}$ consisting of the coordinates of the elements of $R g / R f$ in the basis $\left\{1, t, \ldots, t^{n-1}\right\}$ is called a cyclic ( $f, \sigma, \delta$ )-code.

## Theorem

Let $g(t):=\sum_{i=0}^{r} g_{i} t^{i} \in R$ be a monic right divisor of $f(t)$.
(a) The code corresponding to $R g / R f$ is a free left A-module of dimension $n-r$ where $\operatorname{deg}(f)=n$ and $\operatorname{deg}(g)=r$.
(b) If $v:=\left(a_{0}, a_{1}, \ldots, a_{n-1}\right) \in C$ then $T_{f}(v) \in C$.
(c) The rows of the matrix generating the code $C$ are given by

$$
\left(T_{f}\right)^{k}\left(g_{0}, g_{1}, \ldots, g_{r}, 0, \ldots, 0\right), \quad \text { for } 0 \leq k \leq n-r-1
$$

## Example

Consider $f(t)=t^{5}-1 \in R=\frac{\mathbb{F}_{5}[X]}{X^{5}-1}\left[t ; I d ., \frac{d}{d X}\right]$. In this case $f(x)=f\left(x+x^{4}\right)=0\left(\right.$ with $\left.x=X+\left(X^{5}-1\right)\right)$ and $g(t)=\left[t-x, t-\left(x+x^{4}\right)\right]_{\prime}=t^{2}-2 x t+x^{2}-1$. The generating matrix of the code corresponding to the module $R g / R f$ is given by:

$$
\left(\begin{array}{ccccc}
x^{2}-1 & -2 x & 1 & 0 & 0 \\
2 x & x^{2}+2 & -2 x & 1 & 0 \\
2 & 4 x & x^{2} & -2 x & 1
\end{array}\right)
$$

## Lemma

$f(t), p(t), q(t)=\sum_{i=0}^{n-1} \in R=A[t ; \sigma, \delta]$ such that
$\operatorname{deg}(q(t))<\operatorname{deg}(f(t))=n$. Then
$p(t) q(t) \in R f(t) \Leftrightarrow p\left(T_{f}\right)\left(q_{0}, \ldots, q_{n-1}\right)=(0, \ldots, 0)$

## Theorem

Let $f, g, h, h^{\prime} \in R$ monic such that $f=g h=h^{\prime} g$ and let $C$ denote the code corresponding to the cyclic module $R g / R f$. Then the following statements are equivalent:
(i) $\left(c_{0}, \ldots, c_{n-1}\right) \in C$,
(ii) $\left(\sum_{i=0}^{n-1} c_{i} t^{i}\right) h(t) \in R f$,
(iii) $\sum_{i=0}^{n-1} c_{i} T_{f}^{i}(\underline{h})=\underline{0}$,

## Definition

For a left (resp. right) linear code $C \subseteq A^{n}$, we say that a matrix $H$ is a control matrix if $C=\operatorname{lann}(H)($ resp. $C=\operatorname{rann}(H)$ ).

## Corollary

$f, g, h, h^{\prime} \in R=A[t ; \sigma, \delta]$ monic such that $f=g h=h^{\prime} g$. Then $H={ }^{t}\left(\underline{h}, T_{f}(\underline{h}), \ldots, T_{f}^{\operatorname{deg}(f)-1}(\underline{h})\right)$ is a control matrix of the code corresponding to $R g / R f$.

## Corollary

$f, g, h, h^{\prime} \in R=A[t ; \sigma, \delta]$ monic such that $f=g h=h^{\prime} g$. Then $H={ }^{t}\left(\underline{h}, T_{f}(\underline{h}), \ldots, T_{f}^{\operatorname{deg}(f)-1}(\underline{h})\right)$ is a control matrix of the code corresponding to $R g / R f$.

## Example

$f(t)=t^{5}-1=g(t) h(t)=h(t) g(t) \in R:=\mathbb{F}_{5}[x] /\left(x^{5}-1\right)\left[t ; \frac{d}{d x}\right]$, with $h(t)=t^{3}+2 x t^{2}+\left(3 x^{2}+2\right) t+\left(4 x^{3}+3 x\right)$ and $g(t):=t^{2}-2 x t+x^{2}-1$. C corresponding to $\operatorname{Rg}(t) /\left(t^{5}-1\right)$.

$$
H=\left(\begin{array}{ccccc}
4 x^{3}+3 x & 3 x^{2}+2 & 2 x & 1 & 0 \\
2 x^{2}+3 & 4 x^{3}+4 & 3 x^{2}+4 & 2 x & 1 \\
4 x+1 & 4 x^{2}+2 & 4 x^{3} & 3 x^{2}+1 & 2 x \\
2 x+4 & 2 x+1 & x^{2}+2 & 4 x^{3}+6 x & 3 x^{2}+3 \\
3 x^{2} & 2 x+1 & 4 x+1 & 3 x^{2}+3 & 4 x^{3}+2 x
\end{array}\right)
$$

## $(\sigma, \delta) \mathrm{W}$ codes

## Definitions

a) $f(t) \in R=A[t ; \sigma, \delta]$ is a $W$-polynomial if $f(t)$ is monic and there exist elements $a_{1}, \ldots, a_{n} \in A$ such that $R f(t)=\cap_{i=0}^{i=n} R\left(t-a_{i}\right)$.

## $(\sigma, \delta) \mathrm{W}$ codes

## Definitions

a) $f(t) \in R=A[t ; \sigma, \delta]$ is a $W$-polynomial if $f(t)$ is monic and there exist elements $a_{1}, \ldots, a_{n} \in A$ such that $R f(t)=\cap_{i=0}^{i=n} R\left(t-a_{i}\right)$.
b) The $n \times r$ generalized Vandermonde matrix defined by $a_{1}, \ldots, a_{r}$ is given by:

$$
V_{n}\left(a_{1}, \ldots, a_{r}\right)=\left(\begin{array}{cccc}
1 & 1 & \ldots & 1 \\
a_{1} & a_{2} & \ldots & a_{r} \\
\ldots & \ldots & \ldots & \ldots \\
N_{n-1}\left(a_{1}\right) & N_{n-1}\left(a_{2}\right) & \ldots & N_{n-1}\left(a_{r}\right)
\end{array}\right)
$$

The Wedderburn polynomials play the role of separable polynomials.

## Proposition

Let $f(t), g(t) \in R=A[t ; \sigma, \delta]$ be monic polynomials of degree $n$ and $r$ respectively. Suppose that $g(t)$ is a Wedderburn polynomial with $f(t) \in \operatorname{Rg}(t)$ and let $C$ be the $(\sigma, \delta)-W$-code of length $n$ corresponding to the left cyclic $R$-module $R g(t) / R f(t)$. Let $a_{1}, \ldots, a_{r} \in A$ be such that $R g(t)=\bigcap_{i=0}^{r} R\left(t-a_{i}\right)$. Then $\left(c_{0}, c_{1}, \ldots, c_{n-1}\right) \in C$ if and only if
$\left(c_{0}, c_{1}, \ldots, c_{n-1}\right) V_{n}\left(a_{1}, \ldots, a_{r}\right)=(0, \ldots, 0)$.

## Thanks

## Thank you for your kind attention and ...

## Thanks

## Thank you for your kind attention and ... very mild winter

